

**List 3 Applied Logic**  
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1. Show that a set of *LTL*-tautologies is closed under substitutions, Modus Ponens rule and generalization rule. Is it closed under the following rules:

$$\text{a) } \frac{\varphi}{\diamond\varphi}, \quad \text{b) } \frac{\varphi}{\bigcirc\varphi}, \quad \text{c) } \frac{\varphi}{\varphi U \psi}?$$

2. Prove that the following sentences are *LTL*-tautologies:

$$\begin{array}{ll} \text{a) } \bigcirc\diamond p \leftrightarrow \diamond\bigcirc p, & \text{d) } \diamond\diamond p \leftrightarrow ((q \vee \neg q)U(\diamond p)), \\ \text{b) } (\bigcirc p)U(\bigcirc q) \leftrightarrow \bigcirc(pUq), & \text{e) } (pUq) \wedge \neg q \rightarrow p, \\ \text{c) } (pUq) \vee (rUq) \rightarrow (p \vee qUq), & \text{f) } (\diamond p \wedge \diamond q) \rightarrow \diamond((p \wedge \diamond q) \vee (q \wedge \diamond p)). \end{array}$$

3. Find an equivalent (shortest) version of a sentence:

$$\text{a) } \perp U p, \quad \text{b) } pU \perp, \quad \text{c) } \neg(\perp U \neg p).$$

4. Decide if the sentence is *LTL*-tautology

$$\begin{array}{l} \text{a) } (pUq) \vee (rUq) \leftrightarrow (p \vee qUq), \\ \text{b) } \bigcirc(pUq) \rightarrow \diamond p \wedge \diamond q, \\ \text{c) } (\diamond p \wedge \diamond q) \leftrightarrow \diamond((p \wedge \diamond q) \vee (q \wedge \diamond p)). \end{array}$$

5. Consider a sequence  $(\varphi_n)_{n \in \mathbb{N}}$  defined by the formula:  $\varphi_0 = p$ ,  $\varphi_{n+1} = \bigcirc\varphi_n$  for  $n \in \mathbb{N}$ .

- a) Show that  $\models_{LTL} \varphi_n \leftrightarrow \varphi_m$  if and only if  $n = m$ .  
b) Is the above theorem true after replacing  $\bigcirc$  by  $\diamond$ ?  $\square$ ?

6. Consider Kripke's frame  $(\mathbb{Z}, \leq)$ . We can define connectives  $\bigcirc$  and  $U$ . Let

$$\mathcal{N} = \{(\mathbb{N}, \leq, \pi) : \pi \text{ is a valuation}\},$$

$$\mathcal{Z} = \{(\mathbb{Z}, \leq, \pi) : \pi \text{ is a valuation}\},$$

For a class  $\mathcal{M}$  of Kripke's models set

$$Th_{LTL}(\mathcal{M}) = \{\varphi \in \mathcal{L}_{LTL}(\mathcal{P}) : (\forall M \in \mathcal{M}) M \models \varphi\}.$$

Is it true that  $Th_{LTL}(\mathcal{N}) = Th_{LTL}(\mathcal{Z})$ ?

7. Formulate Principle of Mathematical Induction in *LTL*.

8. Write in *LTL*:

- a)  $p$  will be true exactly twice,  
b)  $q$  will be true but  $p$  will be true earlier,  
c)  $p$  and  $q$  will not be true at the same time.